



Quantum Communication Cost of Preparing Multipartite Entanglement

Caroline Kruszynska, Simon Anders, Wolfgang Dür, Hans-J. Briegel

Institut für Theoretische Physik, Universität Innsbruck, Austria
Institut für Quantenoptik und Quanteninformation der ÖAW, Innsbruck, Austria

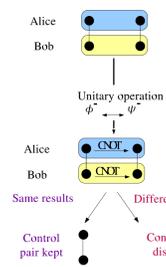


Abstract

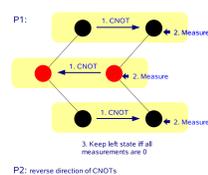
We study the preparation and distribution of high-fidelity multipartite entangled states. There are several possibilities to do that, and which is most efficient depends on the target fidelity one wishes to achieve and on the quality of transmission channel and local operations at one's disposal. We show how to choose the optimal strategy for a variety of settings.

Entanglement Purification Protocols

Bipartite purification
(Oxford protocol)



Multipartite purification
(DAB03; ADB04)



Noise Model and Monte Carlo Simulation

Every two-qubit operation is followed by the noise map $\rho \mapsto (1 - p_1)\rho + p_1 \frac{1}{2}$ acting on the affected qubits. A transmission through the channel is modelled as $\rho \mapsto (1 - p)\rho + p \frac{1}{2}$.

To simulate the effect of this noise, a Monte Carlo technique is employed:

After each operation, the simulator may add an error event. This happens with probability p_1 and is –with equal probabilities– one of the operators in $\mathcal{P} = \{1, \sigma_x, \sigma_y, \sigma_z\}$.

At the end, the remaining states $|\psi\rangle$ are checked by undoing the graph preparation operations and measuring: If

$$\prod_{i \in V} H_i \prod_{\{i,j\} \in E} \Lambda_{Z_{ij}} |\psi\rangle = |0\rangle^{\otimes N},$$

the state is counted as “good”, otherwise “bad” (i. e. erroneous). After summing up counts over many runs, we estimate

$$\text{yield} = \frac{\text{good} + \text{bad}}{\text{initial}}, \quad \text{fidelity} = \frac{\text{good}}{\text{good} + \text{bad}}$$

The Graph State Simulator

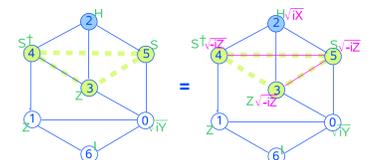
Entanglement purification protocols use only Clifford gates. According to the Gottesman-Knill theorem, they can hence be simulated efficiently on a classical computer.

We have developed such a simulator (AnBr05). Different from the program of Aaronson and Gottesman (AaGo04), it does not directly rely on the stabilizer formalism, but represents a state as

- a graph (given as adjacency list), representing a graph state
- a list of local Clifford operators, which transform this state into the actual one.

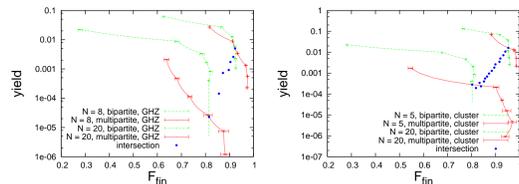
This leads to significantly improved performance.

Measurements are simulated using the rules of (HEB03). Single qubit gates just change the local Clifford operators, and the two-qubit phase gate changes the graph. For the latter, a clever use of the “LU rule” (HEB03; NDM03) is essential.



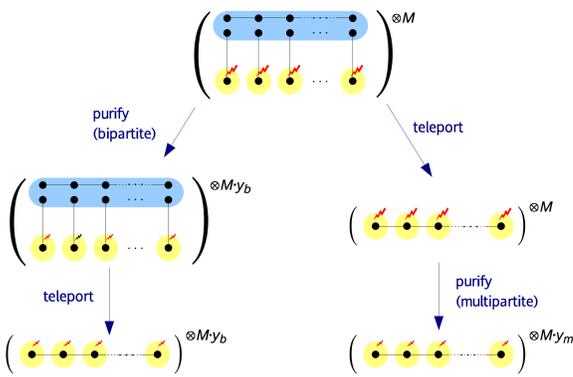
Applicability Range of Both Strategies

Examples for $q = 0.9$ (i. e. 10% channel noise) and $q_l = 0.99$ (i. e. 1% local noise) — preliminary data, unverified



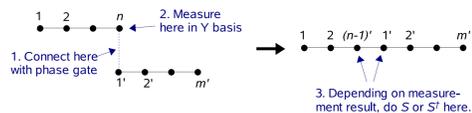
For high target fidelities, multipartite purification gives better yield than bipartite one. The blue disks mark the fidelity/yield values where this break-even occurs for different numbers N of qubits.

Two Strategies to Distribute High Fidelity Entanglement

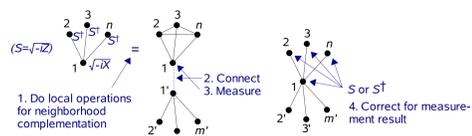


Connecting graph states

It is easy to connect an n -qubits 1D cluster state and an m -qubits cluster state to an $(n + m - 1)$ -qubits cluster state:



To connect GHZ states, we need an additional step:



Analytical recurrence formulae

We could obtain analytical recurrence formulae to describe the multipartite purification protocol for GHZ states for a simplified noise model. Alice's qubit only suffers bit-flipping noise, the other qubits only phase-flipping noise. Then, the state keeps the diagonal, partly permutation-invariant form

$$r_0 |0\rangle \langle 0| + r_1 \sum_{i_1=2}^N |0 \dots 010 \dots\rangle \langle 0 \dots 010 \dots| + r_2 \sum_{\substack{i_1, i_2=2 \\ i_1 < i_2}}^N |0 \dots 1 \dots 1 \dots\rangle \langle \dots | \dots + r_{N-1} |011 \dots 1\rangle \langle \dots|,$$

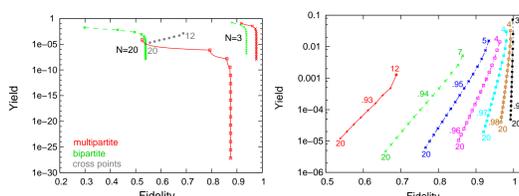
and the effect of noise can be described by a recurrence formula for the coefficients r_i :

$$r'_m = \sum_{l=0}^{N-1} \lambda_{ml} r_l, \\ \lambda_{ml} = \sum_{s=0}^l \binom{N-1}{m} \binom{l}{s} \binom{N-1-l}{m-l+s} (1-q)^{m-l+2s} q^{N-1-m+l-2s}$$

The subsequent purification step can be written this way:

$$r''_m = q r'_m + (1-q) r'_{N+1-m}$$

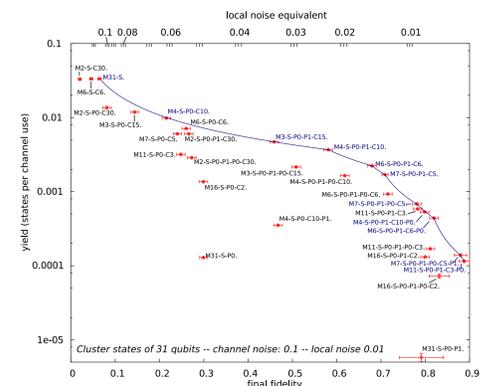
We can use this formulae to study the protocols, and to check our numerics (switched to the simplified noise model) against it.



Results of the analytical model for GHZ states, 10% channel noise, and 7% local noise (left), or 7% to 1% local noise (right).

Intermediate Strategies

The optimal strategy for purification depends on a trade-off between yield and fidelity. Often, it is of advantage to first produce medium-sized states, purify them, connect, and repurify them.



1D cluster states of 31 qubits; channel noise 10%, local noise 1% (preliminary data, unverified)

Legend for the “instruction string”: e. g. “M4-S-P0-C4-P1” means Make 4-qubit states, Send them through the channel, Purify once (direction 0), Connect 4 states to one, Purify again (direction 1).

References

- [AaGo04] Scott Aaronson, Daniel Gottesman: *Improved Simulation of Stabilizer Circuits*, Phys. Rev. A **70** (2004) 052328
- [AnBr05] Simon Anders, Hans J. Briegel: *Fast simulation of stabilizer circuits using a graph state representation* quant-ph/0504117
- [ADB04] Hans Aschauer, Wolfgang Dür and Hans J. Briegel: *Multipartite entanglement purification for two-colorable graph states*, Phys. Rev. A **71** (2005) 012319
- [DAB03] Wolfgang Dür, Hans Aschauer, Hans-J. Briegel: *Multipartite entanglement purification for graph states*, Phys. Rev. Lett. **91** (2003) 107903
- [HEB03] Marc Hein, Jens Eisert, Hans J. Briegel: *Multiparty entanglement in graph states*, Phys. Rev. A **69** (2004) 062311
- [NDM03] Maarten Van den Nest, Jeroen Dehaene, Bart De Moor: *Graphical description of the action of local Clifford transformations on graph states*, Phys. Rev. A **69** (2004) 022316

Acknowledgements

This work was supported in part by the Deutsche Forschungsgemeinschaft (DFG), the Austrian Science Foundation (FWF), the European Union (IST-2001-38877, -39227, OLAQUI, SCALA). W. D. is supported by project APART of the FWF. The numerical calculations has been carried out on the Linux compute cluster of the University of Innsbruck's Konsortium Hochleistungsrechnen.

Comparison of analytical and numerical calculations for the simplified noise model:

